Advanced MPI Parallel File I/O





EURO

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Acknowledgments



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Chap.13 Parallel File I/O

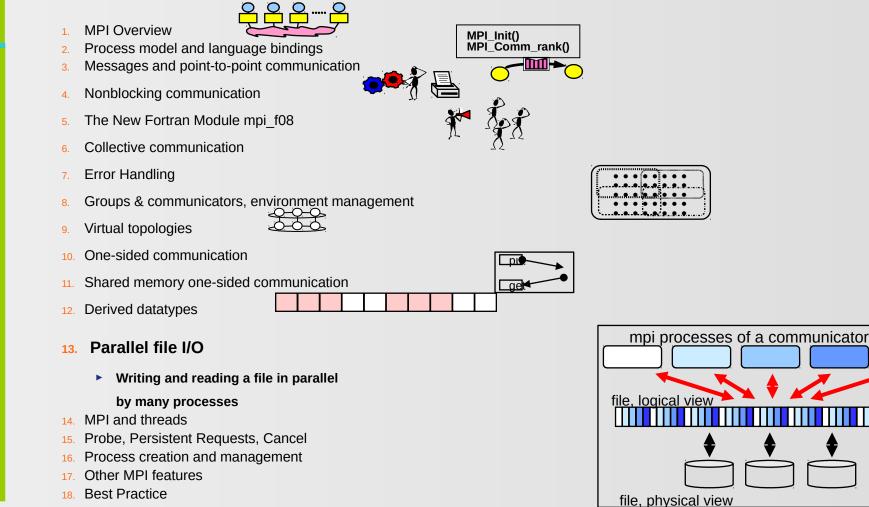


scope of

MPI-I/O

addressed

only by hints



Parallel File I/C

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Motivation, I.



- Many parallel applications need
 - coordinated parallel access to a file by a group of processes
 - simultaneous access
 - all processes may read/write many (small) non-contiguous pieces of the file,
 i.e. the data may be distributed amongst the processes according to a partitioning scheme
 - all processes may read the same data
- Efficient collective I/O based on
 - fast physical I/O by several processors, e.g. striped
 - distributing (small) pieces by fast message passing



Motivation, II.

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- Analogy: writing / reading a file is like sending/receiving a message
- Handling parallel I/O needs
 - handling groups of processes
 - collective operations -> communicators
 - nonblocking operations
 to overlap computation & I/O
 - non-contiguous access

- MPI topologies and groups
 - file handle defined like
 - MPI_I..., MPI_Wait, ... new **split** collective interface
 - MPI derived datatypes



MPI-I/O Features



- Provides a high-level interface to support
 - data file partitioning among processes
 - transfer global data between memory and files (collective I/O)
 - asynchronous transfers
 - strided access
- MPI derived datatypes used to specify common data access patterns for maximum flexibility and expressiveness



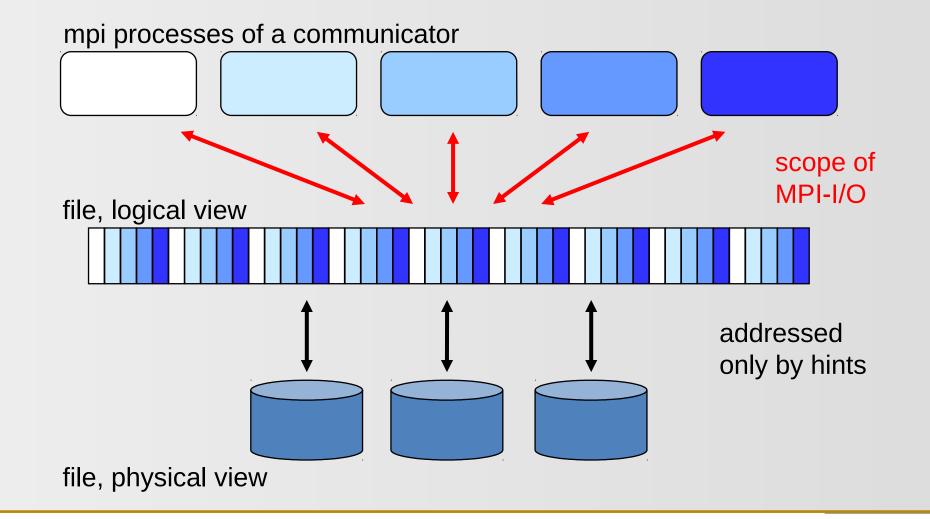
MPI-I/O Principles

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- MPI file contains elements of a single MPI datatype (etype)
- partitioning the file among processes with an access template (filetype)
- all file accesses transfer to/from a contiguous or non-contiguous user buffer (MPI datatype)
- nonblocking / blocking and collective / individual read / write routines
- individual and shared file pointers, explicit offsets
- binary I/O
- automatic data conversion in heterog. systems
- file interoperability with external representation



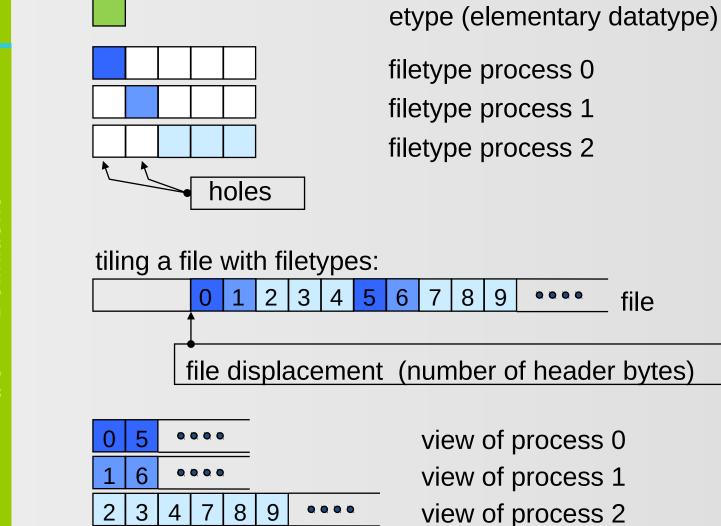




Definitions

file





Comments on Definitions



- an ordered collection of typed data items
- etypes is the unit of data access and positioning / offsets
 - can be any basic or derived datatype (with non-negative, monotonically non-decreasing, non-absolute displacem.)
 - generally contiguous, but need not be
 - typically same at all processes
- filetypes the basis for partitioning a file among processes
 - defines a template for accessing the file
 - different at each process
 - the etype or derived from etype (displacements: non-negative, monoton. non-decreasing, non-abs., <u>multiples of etype extent</u>)
- view each process has its own view, defined by: a displacement, an etype, and a filetype.
 - The filetype is repeated, starting at **displacement**
- offset position relative to current view, in units of etype

file

Opening an MPI File



- MPI_FILE_OPEN is collective over comm
- filename's namespace is implementation-dependent!
- filename must reference the same file on all processes
- process-local files can be opened by passing MPI_COMM_SELF as comm
- returns a file handle fh

[represents the file, the process group of **comm**, and the current view]

MPI_FILE_OPEN(comm, filename, amode, info, *fh*)



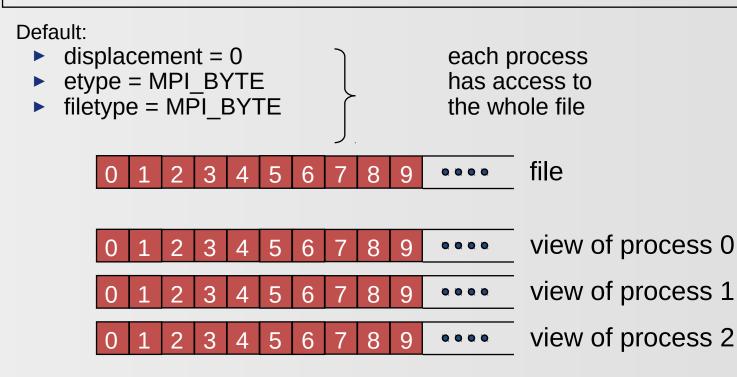
C/C++ language bindings – see MPI Standard



Default View



MPI_FILE_OPEN(comm, filename, amode, info, *fh*)



- Sequence of MPI_BYTE matches with any datatype (see MPI-3.0, Section 13.6.5 on page 544 / MPI-3.1, Section 13.6.6 on page 549)
- Binary I/O (no ASCII text I/O)



Access Modes



- same value of amode on all processes in MPI_FILE_OPEN
- Bit vector OR of integer constants (Fortran 77: +)
 - MPI_MODE_RDONLY read only
 - MPI_MODE_RDWR reading and writing
 - MPI_MODE_WRONLY write only
 - MPI_MODE_CREATE create if file doesn't exist
 - MPI_MODE_EXCL error creating a file that exists
 - MPI_MODE_DELETE_ON_CLOSE delete on close
 - MPI_MODE_UNIQUE_OPEN file not opened concurrently
 - MPI_MODE_SEQUENTIAL file only accessed sequentially: mandatory for sequential stream files (pipes, tapes, ...)
 - MPI_MODE_APPEND all file pointers set to end of file [caution: reset to zero by any subsequent MPI_FILE_SET_VIEW]



File Info: Reserved Hints



- Argument in MPI_FILE_OPEN, MPI_FILE_SET_VIEW, MPI_FILE_SET_INFO
- reserved key values:
 - collective buffering
 - "collective_buffering": specifies whether the application may benefit from collective buffering
 - "cb_block_size": data access in chunks of this size
 - "cb_buffer_size": on each node, usually a multiple of block size
 - "cb_nodes": number of nodes used for collective buffering
 - disk striping (only relevant in MPI_FILE_OPEN)
 - "striping_factor": number of I/O devices used for striping
 - "striping_unit": length of a chunk on a device (in bytes)
- MPI_INFO_NULL may be passed



Closing and Deleting a File



Close: collective

MPI_FILE_CLOSE(fh)

- Delete:
 - automatically by MPI_FILE_CLOSE

if amode=MPI_DELETE_ON_CLOSE | ...

was specified in MPI_FILE_OPEN

deleting a file that is not currently opened:

MPI_FILE_DELETE(filename, info)

[same implementation-dependent rules as in MPI_FILE_OPEN]

Writing with Explicit Offsets



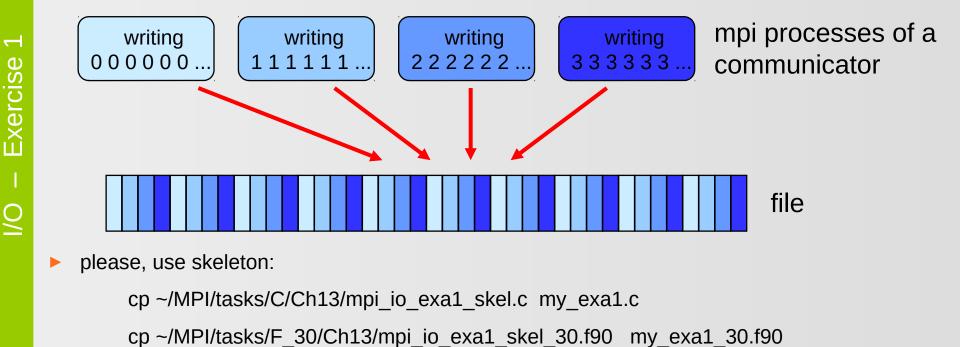
MPI_FILE_WRITE_AT(fh,offset,buf,count,datatype,status)

- writes count elements of datatype from memory buf to the file
- starting offset * units of etype from begin of view
- the elements are stored into the locations of the current view
- the sequence of basic datatypes of datatype (= signature of datatype) must match contiguous copies of the etype of the current view

Explicit Offsets - WRITE / MPI–IO Exercise 1: Four processes write a file in parallel



- each process should write its rank (as one character) ten times to the offsets = my rank + i * size of MPI COMM WORLD, i=0..9
- Each process uses the default view



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0

During the Exercise





Please stay here in the main room while you do this exercise

And have fun with this short exercise



Please do not look at the solution before you finished this exercise,

otherwise,

90% of your learning outcome may be lost

As soon as you finished the exercise,

please go to your breakout room

and continue your discussions with your fellow learners:

Who of you uses NetCDF or HDF5?

As far as I know, both use MPI-I/O.



Outline – Block 2



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File Views



- Provides a visible and accessible set of data from an open file
- A separate view of the file is seen by each process through <u>triple :=</u> (displacement, etype, filetype)
- User can change a view during the execution of the program <u>but collective operation</u>
- A linear byte stream, <u>represented by the triple</u>
 - (0, MPI_BYTE, MPI_BYTE), is the default view



Set/Get File View



Set view

- changes the process's view of the data
- local and shared file pointers are reset to zero
- <u>collective</u> operation
- etype and filetype must be committed
- datarep argument is a string that specifies the format in which data is written to a file:
 - "native", "internal", "external32", or user-defined
- same etype extent and same datarep on all processes
- Get view
 - returns the process's view of the data

MPI_FILE_SET_VIEW(fh, disp, etype, filetype, datarep, info) MPI_FILE_GET_VIEW(fh, disp, etype, filetype, datarep)

Data Representation, I.



"native"

- data stored in file identical to memory
- on homogeneous systems no loss in precision or I/O performance due to type conversions
- on heterogeneous systems loss of interoperability
- no guarantee that MPI files accessible from C/Fortran

"internal"

- data stored in implementation specific format
- can be used with homogeneous or heterogeneous environments
- implementation will perform type conversions if necessary
- no guarantee that MPI files accessible from C/Fortran

Data Representation, II.



"external32"

- follows standardized representation (IEEE)
- all input/output operations are converted from/to the "external32" representation
- files can be exported/imported between different MPI environments
- due to type conversions from (to) native to (from) "external32" data precision and I/O performance may be lost
- "internal" may be implemented as equal to "external32"
- can be read/written also by non-MPI programs
- user-defined

No information about the default,

i.e., datarep without MPI_File_set_view() is not defined

Fileview examples with SUBARRAY and DARRAY



Task

- reading a global matrix from a file
- storing a subarray into a local array on each process
- according to a given distribution scheme



Example with Subarray, I.



- 2-dimensional distribution scheme: (BLOCK, BLOCK)
- garray on the file 20x30:
 - Contiguous indices is language dependent:
 - in Fortran: (1,1), (2,1), (3,1), ..., (1,10), (2,20), (3,10), ..., (20,30)
 - ▶ in C/C++: [0][0], [0][1], [0][2], ..., [10][0], [10][1], [10][2], ..., [19][29]
- larray = local array in each MPI process
 - = subarray of the global array
- same ordering on file (garray) and in memory (larray)

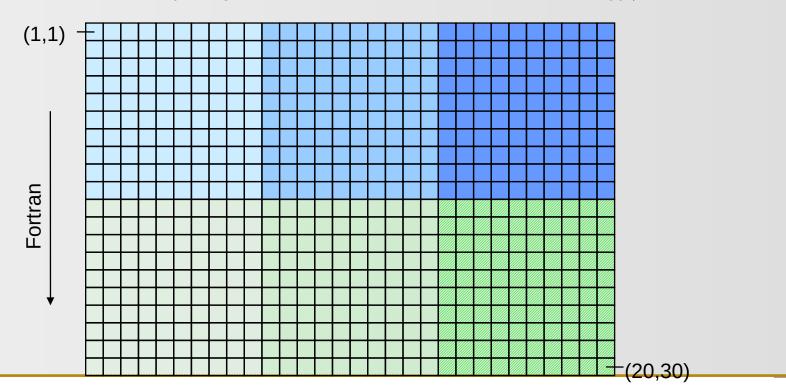
Example with Subarray, II. — Distribution



Process topology: 2x3

- global array on the file: 20x30
- distributed on local arrays in each processor: 10x10

C / C++ (contiguous indices on the file and in the memory)



Example with Subarray, III. — Reading the file

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```
! these HPF-like comment lines !
!!!! real garray(20,30)
!!!! PROCESSORS procs(2, 3)
                                                 ! explain the data distribution
   DISTRIBUTE garray(BLOCK, BLOCK) onto procs ! used in this MPI program
1111
   real larray(10,10); integer (kind=MPI_OFFSET_KIND) disp,offset; disp=0; offset=0
   ndims=2; psizes(1)=2; period(1)=.false.; psizes(2)=3; period(2)=.false.
   call MPI_CART_CREATE(MPI_COMM_WORLD, ndims, psizes, period,
                                                      .TRUE., comm, ierror)
   call MPI COMM RANK(comm, rank, ierror)
   call MPI_CART_COORDS(comm, rank, ndims, coords, ierror)
   gsizes(1)=20; lsizes(1)=10; starts(1)=coords(1)*lsizes(1)
   gsizes(2)=30; lsizes(2)=10; starts(2)=coords(2)*lsizes(2)
   call MPI_TYPE_CREATE_SUBARRAY(ndims, gsizes, lsizes, starts,
                   MPI_ORDER_FORTRAN, MPI_REAL, subarray_type, ierror)
   call MPI TYPE COMMIT(subarray type, ierror)
   call MPI FILE OPEN(comm, 'exa subarray testfile', MPI MODE CREATE +
                         MPI_MODE_RDWR, MPI_INFO_NULL, fh, ierror)
   call MPI_FILE_SET_VIEW (fh, disp, MPI_REAL, subarray_type, 'native',
                              MPI INFO NULL, ierror)
   call MPI FILE READ AT ALL(fh, offset, larray, lsizes(1)*lsizes(2), MPI_REAL,
                                  status, ierror)
```

Example with Subarray, IV.



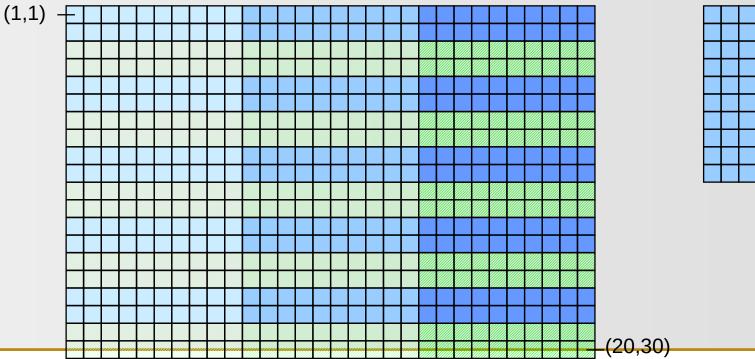
- All MPI coordinates and indices start with 0, even in Fortran, i.e. with MPI_ORDER_FORTRAN
- MPI indices (here starts) may differ () from Fortran indices
- Block distribution on 2*3 processes:

rank = 0 coords = (0, 0) starts = (0, 0) garray(1:10, 1:10) = larray (1:10, 1:10)		
rank = 3 coords = $(1, 0)$ starts = $(10, 0)$ garray(11:20, 1:10) = larray (1:10, 1:10)	rank = 4 coords = (1, 1) starts = (10, 10) garray(11:20, 11:20) = larray (1:10, 1:10)	rank = 5 coords = $(1, 2)$ starts = $(10, 20)$
		,

Example with Darray, I.

Distribution scheme: (CYCLIC(2), BLOCK)

- Cyclic distribution in first dimension with strips of length 2
- Block distribution in second dimension
- distribution of global garray onto the larray in each of the 2x3 processes
- garray on the file: e.g., larray on process (0,1):





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Example with Darray, II.



```
!!!!! real garray(20,30)
                                               ! these HPF-like comment lines !
IIII PROCESSORS procs(2, 3)
                                                  explain the data distribution!
ITE DISTRIBUTE garray(CYCLIC(2), BLOCK) onto proce I used in this MPI program!
   real larray(10,10); integer (kind=MPI_OFFSET_KIND) disp, offset; disp=0; offset=0
   call MPI COMM SIZE(comm, size, ierror)
    ndims=2; psizes(1)=2; period(1)=.false.; psizes(2)=3; period(2)=.false.
   call MPI_CART_CREATE(MPI_COMM_WORLD, ndims, psizes, period,
                                                      .TRUE., comm, ierror)
   call MPI COMM RANK(comm, rank, ierror)
   call MPI_CART_COORDS(comm, rank, ndims, coords, ierror)
   gsizes(1)=20; distribs(1)= MPI_DISTRIBUTE_CYCLIC; dargs(1)=2
   gsizes(2)=30; distribs(2)= MPI_DISTRIBUTE_BLOCK; dargs(2)=
                                            MPI DISTRIBUTE DFLT DARG
   call MPI_TYPE_CREATE_DARRAY(size, rank, ndims, gsizes, distribs, dargs,
             psizes, MPI_ORDER_FORTRAN, MPI_REAL, darray_type, ierror)
   call MPI_TYPE_COMMIT(darray_type , ierror)
   call MPI FILE OPEN(comm, 'exa subarray testfile', MPI MODE CREATE +
                         MPI_MODE_RDWR, MPI_INFO_NULL, fh, ierror)
   call MPI_FILE_SET_VIEW (fh, disp, MPI_REAL, darray_type, 'native',
                              MPI_INFO_NULL, ierror)
   call MPI FILE READ AT ALL(fh, offset, larray, 10*10, MPI REAL, istatus, ierror)
```

Example with Darray, III.



- Cyclic distribution in first dimension with strips of length 2
- Block distribution in second dimension
- Processes' tasks:

rank = 0 coords = (0, 0) $\begin{bmatrix} 1: 2\\5: 6\\9:10\\13:14\\17:18\\\end{bmatrix}$ = larray (1:10, 1:10)	rank = 1 coords = (0, 1) $\begin{bmatrix} 1: 2\\ 5: 6\\ 9:10\\ 13:14\\ 17:18\\ 17:18\\ 11:10\\ 1$	rank = 2 coords = (0, 2) $\begin{bmatrix} 1: 2\\ 5: 6\\ 9:10\\ 13:14\\ 17:18\\ \end{bmatrix}$ = larray (1:10, 1:10)
rank = 3 coords = $(1, 0)$ 3: 4 7: 8 11:12 1:10) 15:16 19:20 = larray (1:10, 1:10)	rank = 4 coords = $(1, 1)$ $\begin{bmatrix} 3: 4\\ 7: 8\\ 7: 8\\ 11:12\\ 15:16\\ 19:20\\ \end{bmatrix}$ = larray (-1:10, 1:10)	rank = 5 coords = $(1, 2)$ $\begin{bmatrix} 3: 4\\ 7: 8\\ 11:12\\ 15:16\\ 19:20 \end{bmatrix}$ = larray (-1:10, 1:10)

5 Aspects of Data Access



- Direction: Read / Write
- Positioning [realized via routine names]
 - explicit offset (_AT)
 - individual file pointer (no positional qualifier)
 - shared file pointer (_SHARED or _ORDERÉD) (different names used depending on whether non-collective)
 - (different names used depending on whether non-collective or collective)
- Coordination
 - non-collective
 - collective (_ALL)
- Synchronism
 - blocking
 - nonblocking (I) and split collective (_BEGIN, _END)
- Atomicity, [realized with a separate API: MPI_File_set_atomicity]
 - non-atomic (default)
 - atomic: to achieve sequential consistency for conflicting accesses on same fh in different processes

All Data Access Routines



positioning	synchronism	coordina	ation			
		noncollective	collective	split collective		
explicit	blocking	READ_AT	READ_AT_ALL	READ_AT_ALL_BEGIN		
offsets		WRITE_AT	WRITE_AT_ALL	READ_AT_ALL_END		
	nonblocking	IREAD_AT	IREAD_AT_ALL	WRITE_AT_ALL_BEGIN		
		IWRITE_AT	IWRITE_AT_ALL	WRITE_AT_ALL_END		
individual	blocking	READ	READ_ALL	READ_ALL_BEGIN		
file pointers		WRITE	WRITE_ALL	READ_ALL_END		
	nonblocking	IREAD	IREAD_ALL	WRITE_ALL_BEGIN		
		IWRITE	IWRITE_ALL	WRITE_ALL_END		
shared	blocking	READ_SHARED	READ_ORDERED	READ_ORDERED_BEGIN		
file pointer		WRITE_SHARED	WRITE_ORDERED	READ_ORDERED_END		
	nonblocking	IREAD_SHARED	N/A	WRITE_ ORDERED_ BEGIN		
		IWRITE_SHARED		WRITE_ ORDERED_ END		
Read e.g. MPI_FILE_READ_AT New in MPI-3.1						

Explicit Offsets



e.g. MPI_FILE_READ_AT(fh,offset,*buf*,count,datatype,*status*)

- attempts to read count elements of datatype
- starting offset * units of etype
 from begin of view (= displacement)
- the sequence of basic datatypes of datatype
 (= signature of datatype)
 - (= signature of **datatype**)
 - must match

contiguous copies of the etype of the current view

- EOF can be detected by noting that the amount of data read is less than count
 - i.e. EOF is no error!
 - use MPI_GET_COUNT(status,datatype,recv_count)

Individual File Pointer, I.



e.g. MPI_FILE_READ(fh, *buf*,count,datatype,*status*)

- same as *"Explicit Offsets"*, except:
- the offset is the current value of the individual file pointer of the calling process
- the individual file pointer is updated by

new_fp = old_fp + elements(datatype) * count

i.e. it points to the next etype after the last one that will be accessed (formula is not valid if EOF is reached)

Individual File Pointer, II.



MPI_FILE_SEEK(fh, offset, whence)

- set individual file pointer fp:
 - set fp to offset if whence=MPI_SEEK_SET
 - advance fp by offset if whence=MPI_SEEK_CUR
 - set fp to EOF+offset if whence=MPI_SEEK_EOF

MPI_FILE_GET_POSITION(fh, offset)
MPI_FILE_GET_BYTE_OFFSET(fh, offset, disp)

- to inquire offset
- to convert offset into byte displacement
 - [e.g. for disp argument in a new view]



Copy to your local directory:

cp ~/MPI/tasks/C/Ch13/mpi_io_exa2_skel.c my_exa2.c

cp ~/MPI/tasks/F_30/Ch13/mpi_io_exa2_skel_30.f90 my_exa2_30.f90

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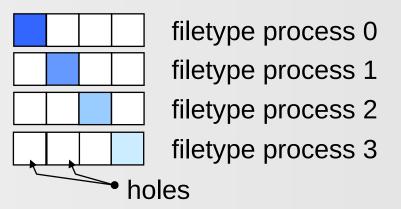
Tasks:

- Each MPI-process of my_exa2 should write one character to a file:
 - process "rank=0" should write an 'a'
 - process "rank=1" should write an 'b'
 - <u>►</u> ...
- Use a 1-dimensional fileview with MPI_TYPE_CREATE_SUBARRAY
- The pattern should be repeated 3 times, i.e., four processes should write: "abcdabcdabcd"
- Please, substitute "_____" in your my_exa2.c / _30.f90
- Compile and run your my_exa2.c / _30.f90

MPI-IO Exercise 2: Using fileviews and individual filepointers, continued



etype = MPI CHARACTER / MPI CHAR



tiling a file with filetypes:

file b c d a b c d a b c d

file displacement = 0 (number of header bytes)

a	a	a	0.0.0.0	view of process 0
b	b	b		view of process 1
С	С	С		view of process 2
d	d	d		view of process 3







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Please stay here in the main room while you do this exercise

And have fun with this short exercise



Please do not look at the solution before you finished this exercise,

otherwise,

90% of your learning outcome may be lost



As soon as you finished the exercise, please go to your breakout room and continue your discussions with your fellow learners:

Ask yourself, whether the datatype is a 1- or higher-dimensional array?

And don't forget that counts are normally elements and not bytes!



And to look at the declaration of the buffer is also helpful

to answer tha last _____ question

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Shared File Pointer, I.



- same view at all processes mandatory!
- the offset is the current, global value of the

shared file pointer of fh

- multiple calls [e.g. by different processes] behave as if the calls were serialized
- non-collective, e.g.

MPI_FILE_READ_SHARED(fh, *buf*, count, datatype, *status*)

collective calls are serialized in the order of the processes' ranks, e.g.:

MPI_FILE_READ_ORDERED(fh,*buf*,count,datatype,*status*)



Shared File Pointer, II.



MPI_FILE_SEEK_SHARED(fh, offset, whence)
MPI_FILE_GET_POSITION_SHARED(fh, offset)
MPI_FILE_GET_BYTE_OFFSET(fh, offset, disp)

same rules as with individual file pointers

Collective Data Access



- Explicit offsets / individual file pointer:
 - same as non-collective calls by all processes "of fh"
 - opportunity for best speed!!!
- shared file pointer:
 - accesses are ordered by the ranks of the processes
 - optimization opportunity:
 - first, locations within the file for all processes can be computed
 - then parallel physical data access by all processes



Application Scenery, I.



Scenery A:

- Task: Each process has to read the whole file
- Solution: MPI_FILE_READ_ALL
 - = collective with individual file pointers,
 - with same view (displacement+etype+filetype)
 - on all processes
 - [internally: striped-reading by several process, only once
 - from disk, then distributing with bcast]
- Scenery B:
 - Task: The file contains a list of tasks,
 - each task requires different compute time
 - Solution: MPI_FILE_READ_SHARED
 - = non-collective with a shared file pointer
 - (same view is necessary for shared file p.)



Application Scenery, II.



Scenery C:

- Task: The file contains a list of tasks,
 - each task requires the same compute time
- Solution: MPI_FILE_READ_ORDERED
 - = collective with a shared file pointer
 - (same view is necessary for shared file p.)
- or: MPI_FILE_READ_ALL
 - = collective with individual file pointers,
 - different views: filetype with
 - MPI_TYPE_CREATE_SUBARRAY(1,nproc,
 - 1, myrank, ..., datatype_of_task, *filetype*)
 - [internally: both may be implemented the same
 - and equally with following scenery D]



Application Scenery, III.



Scenery D:

- Task: The file contains a matrix,
 - block partitioning,
 - each process should get a block
- Solution: generate different filetypes with
 - MPI_TYPE_CREATE_DARRAY,
 - the view on each process represents the block
 - that should be read by this process,
 - MPI_FILE_READ_AT_ALL with offset=0
 - (= collective with explicit offsets)
 - reads the whole matrix collectively
 - [internally: striped-reading of contiguous blocks
 - by several process,
 - then distributed with "alltoall"]

Nonblocking Data Access



Split Collective Non-Blocking

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e.g. MPI_FILE_IREAD(fh, *buf*, count, datatype, *request*) MPI_WAIT(request, *status*) MPI_TEST(request, *flag, status*)

analogous to MPI-1 nonblocking

Split Collective Data Access, I.



- collective operations may be **split** into two parts:
 - start the split collective operation

e.g. MPI_FILE_READ_ALL_BEGIN(fh, *buf*, count, datatype)

complete the operation and return the status

MPI_FILE_READ_ALL_END(fh, buf, status)

Split Collective Data Access, II.



- Rules and Restrictions:
 - the MPI_...BEGIN calls are collective
 - the MPI_...END calls are collective, too
 - only one active (pending) split or regular collective operation per file handle at any time
 - split collective does not match ordinary collective
 - same buf argument in MPI_...BEGIN and ..._END call
- opportunity to overlap file I/O and computation
- but also a valid implementation:
 - does all work within the MPI_...BEGIN routine, passes status in the MPI_...END routine
 - passes arguments from MPI_...BEGIN to MPI_...END, does all work within the MPI ...END routine



Scenery – Split Collective



Scenery A:

- Task:Each process has to read the whole file
- Solution: MPI_FILE_READ_ALL_BEGIN

= collective with individual file pointers,
with same view (displacement+etype+filetype)
on all processes
[internally: starting asynchronous striped-reading
by several process]

- then computing some other initialization,
- MPI_FILE_READ_ALL_END.

[internally: waiting until striped-reading finished, then distributing the data with bcast]

Other File Manipulation Routines



- Pre-allocating space for a file [collective call, may be expensive]
 MPI_FILE_ PREALLOCATE(fh, size)
- Resizing a file [collective call, may speed up first writing on a file] MPI_FILE_SET_SIZE(fh, size)
- Querying file size
 MPI_FILE_GET_SIZE(filename, *size*)
- Querying file parameters
 MPI_FILE_GET_GROUP(fh, group)
 MPI_FILE_GET_AMODE(fh, amode)
- File info object

MPI_FILE_SET_INFO(fh, info) [collective call]

MPI_FILE_GET_INFO(fh, info_used)

Returns a new info object that contains the current setting of **all hints** used by the system related to this open file:

- provided by the application, and
- provided by the system

MPI I/O Error Handling



- File handles have their own error handler
- Default is MPI_ERRORS_RETURN,
 - i.e. non-fatal

[vs message passing: MPI_ERRORS_ARE_FATAL]

Default is associated with MPI_FILE_NULL

[vs message passing: with MPI_COMM_WORLD]

Changing the default, e.g., after MPI_Init:

MPI_File_set_errhandler(MPI_FILE_NULL, MPI_ERRORS_ARE_FATAL);

- CALL MPI_FILE_SET_ERRHANDLER(MPI_FILE_NULL,MPI_ERRORS_ARE_FATAL,*ierr*)
- MPI is undefined after first erroneous MPI call
- but a high quality implementation

will support I/O error handling facilities

Implementation-Restrictions



- ROMIO based MPI libraries:
 - datarep = "internal" and "external32" is still not implemented
 - User-defined data representations are not supported

MPI-I/O: Summary



- Rich functionality provided to support various data representation and access
- MPI I/O routines provide flexibility as well as portability
- Collective I/O routines can improve I/O performance
- Initial implementations of MPI I/O available

(eg, ROMIO from Argonne)

Available nearly on every MPI implementation

MPI–IO Exercise 3: Collective ordered I/O



Copy to your local directory:

cp ~/MPI/tasks/C/Ch13/mpi_io_exa3_skel.c my_exa3.c

cp ~/MPI/tasks/**F_30**/Ch13/mpi_io_exa3_skel_30.f90 **my_exa3_30.f90**

- Tasks:
 - Substitute the write call with individual filepointers
 by a collective write call with shared filepointers
 - Compile and run your my_exa3.c / _30.f90





Please stay here in the main room while you do this exercise

And have fun with this short exercise



Please do not look at the solution before you finished this exercise,

otherwise,

90% of your learning outcome may be lost



As soon as you finished the exercise,

please go to your breakout room

and continue your discussions with your fellow learners:

This exercise is mainly removing all about the fileview. With the shared file pointer and collective writing, this exercise

is a one-line problem, isn't it?

Good luck!

MPI–IO Exercise 4: I/O Benchmark



Use:

MPI/tasks/F_30/Ch13/mpi_io_exa4_30.f90

(my apologies that there is only a Fortran version)

Tasks:

- Compile and execute mpi_io_exa4 on 2, 4 and 8 MPI processes.
- Duplicate "WRITE_ALL & READ_ALL" block

and substitute by non-collective "WRITE & READ".

- Compare collective and non-collective I/O.
- Double the value of gsize and compile and execute again.

Chapter 12-(1), Exercise 1: MPI_TYPE_CONTIGUOUS Scenery, III.

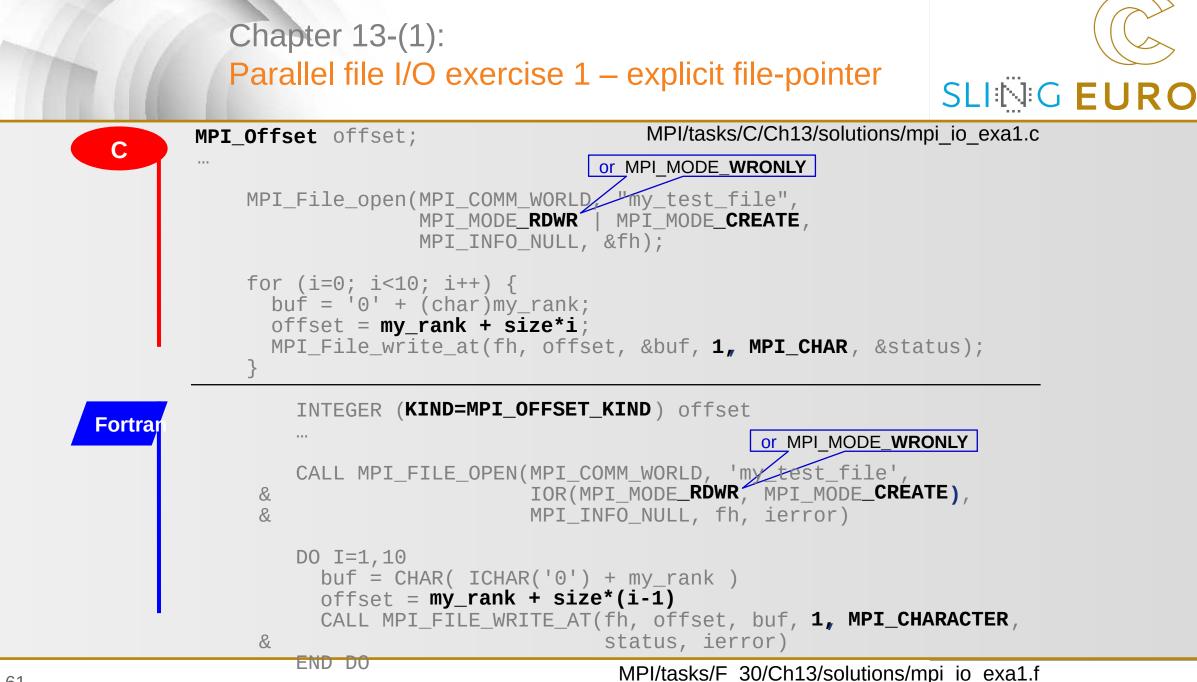


<pre>struct buff{ int i; int j; } snd_buf, rcv_buf, s</pre>	int i; Provided in				
TYPE(MPI_Datatype) ::					
<pre>snd_buf.i = my_rank; for(i = 0; i < size; { MPI_Issend(&snd_buf MPI_Recv (&rcv_buf MPI_Wait(&request, snd_buf = rcv_buf; sum.i += rcv_buf.i; }</pre>	<pre>i,1,send_recv_type,right,17,MPI_COMM_WORLD, &request); i,1,send_recv_type,left, 17,MPI_COMM_WORLD, &status); &status); sum.j += rcv_buf.j;</pre>				
printf ("PE %i: Sum =	* %i and %i \n", my_rank, sum.i, sum.j);				





APPENDIX: Solution to exercises



Chapter 13-(2): Parallel file I/O exercise 2 – with fileview

SLING MPI/tasks/C/Ch13/solutions/mpi io exa2.c MPI_Offsetdisp; С ndims = 1EURC $_{\pm}$ size. array_of_sizes[0] array_of_subsizes[0] = 1; or MPI MODE_WRONLY _ my_rank l array_of_starts[0] MPI_Type_committ_(&fialretaype)); #RDWR CREATE MPI_File_open(..., MPI_MODE MPI MODE , ...); disp = ;or MPI_CHAR MPI_File_set_view(...); for (i=0; i<3; i++) {
 buf = 'a' + (char)my_rank; 1, etype</pre> MPT File write(fh &buf &status): _KIND=MPI_OFFSET_KIND) dMPI/tasks/F 30/Ch13/solutions/mpi io exa2.f INTEGER size Fortrar 1 ndims = = my_rank array_of_sizes(1) or MPI MODE WRONLY array_of_subsizes(1) = array_of_startscommit(filetype, ierror) _RDWR CREATE CALL MPT_TYPE_CREATE_SUBARRAY(...) CALL MPI TYPE CALL MPI_FILE_OPEN(..., IOR(MPI_MODE or MPIPCHARACTER), ...) disp = 1, etype CALL MPI_FILE_SET_VIEW(...) DO I=1,3 buf = CHAR(ICHAR('a') + my_rank) CALL MPI_FILE_WRITE(fh, buf, , , status, ierror) END DO 62

Chapter 13-(3). Parallel file I/O exercise 3 – shared filepointer SLI

EURC MPI/tasks/C/Ch13/solutions/mpi io exa3.c **PI** Datatype etype; С atvne filetvne Offset disp: of_sizes[0] = size; arrav of subsizes[0] = 1;array_of starts[0] = my rank; order = MPX ORDER C; _Type_create_subarray(ndims, array_of_sizes, array of subsizes, array of starts, order, etype, &filetype); Fvpe commit(&filetvpe); 1sp = 0;File_set_view(th, disp, etype, filetype, "native", MPI_INFO_NULL); MPI File open(MPN COMM WORLD, "my test file", MPI MODE RDWR | MPI MODE CREATE, MPI INFO NULL, &fh); for (i=0; i<3; i++) { buf = 'a' + (char)my_rank; MPI_File_write_ordered(fh, &buf, 1, MPI_CHAR, &status); MPI File close(&fh); TYPE(MPI_Datatype) :: etype EILE_OPEN(MPI_COMM_WORLD, 'my_test_File_, & MPI/tasks/F_30/Ch13/solutions/mpi_io_exa3.f Fortran IOR(MPI MODE RDWR, MPI MODE CREATE), & MPI INFO NULL, fh, ierror) DO I=1,3 buf = CHAR(ICHAR('a') + my_rank) CALL MPI_FILE_WRITE_ORDERED(fh, buf, 1, MPI_CHARACTER, status, ierror) END DO CALL MPI_FILE_CLOSE(fh, ierror)



Thanks!





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